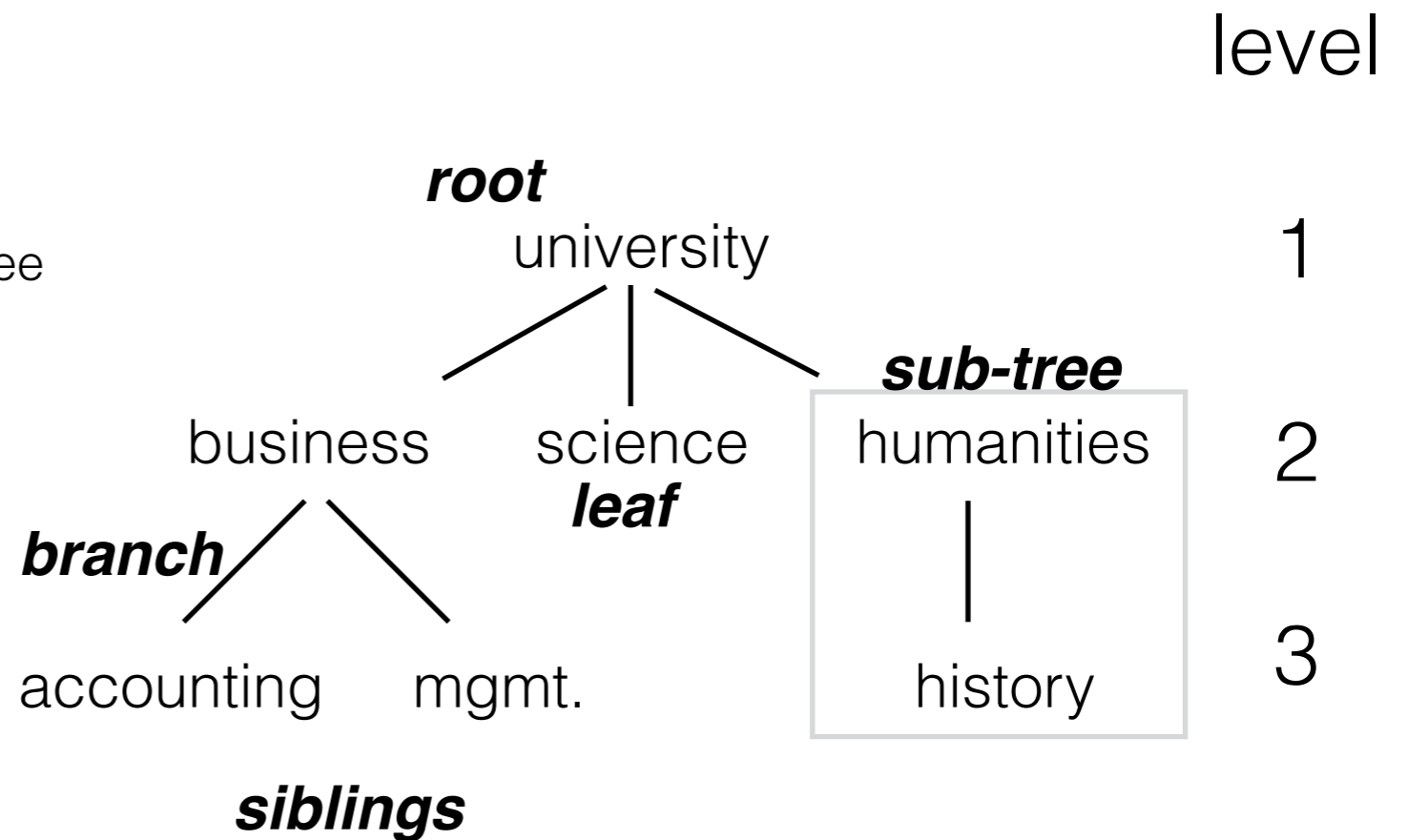


Search Algorithms, Trees, and Graphs Part I

Anton Gerdelan <gerdela@scss.tcd.ie>

General Trees

- **Node** - item in the tree
- **Branch** - link connecting nodes
- **Sub-tree** - part of tree that is also a tree
- **Root** - the node with no **parent**
- **Leaf** - nodes with no **children**
- **Siblings** - children of same parent
- **Levels** - Number of rows
- **Null tree** - nothing in tree
- **Tree** is either
 - a null tree
 - a root with several sub-trees



Typical Tree Operations

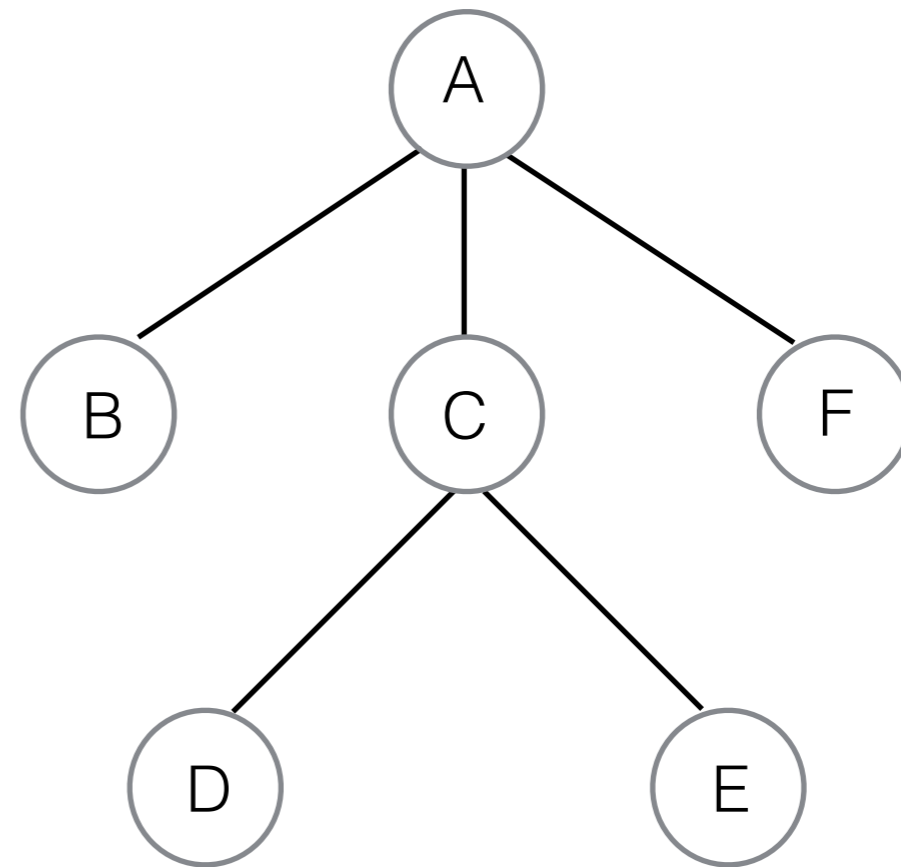
- **data(Node* N)** - return value of data in a node
- **parent(Node* N)** - return pointer to parent
- **leftmost_child(Node* N)**
 - return pointer to furthest left child of N
- **right_sibling(Node* N)**
- **insert_child(Node* N, data)**
 - make a new node with data and make it leftmost child of N
- **is_empty()** - return true if tree empty
- *How many pointers will your tree Node have?*

Searching for Data in a Tree

- An algorithm that visits each node once - **tree traversal** - and compares contents to sought value
- traversal ops: **V** - **visit** / look at node, **L** - **left** sub-tree, **R** - **right** sub-tree
- **Pre-order: V, L to R**
 1. look at data in root
 2. recurse all subtrees from left to right
- **In-order: L, V to R**
 1. recurse leftmost subtree
 2. look at root of current sub-tree
 3. recurse remaining subtrees to rightmost subtree
- **Post-order: L to R, V**
 1. recurse all subtrees from left to right
 2. look at root

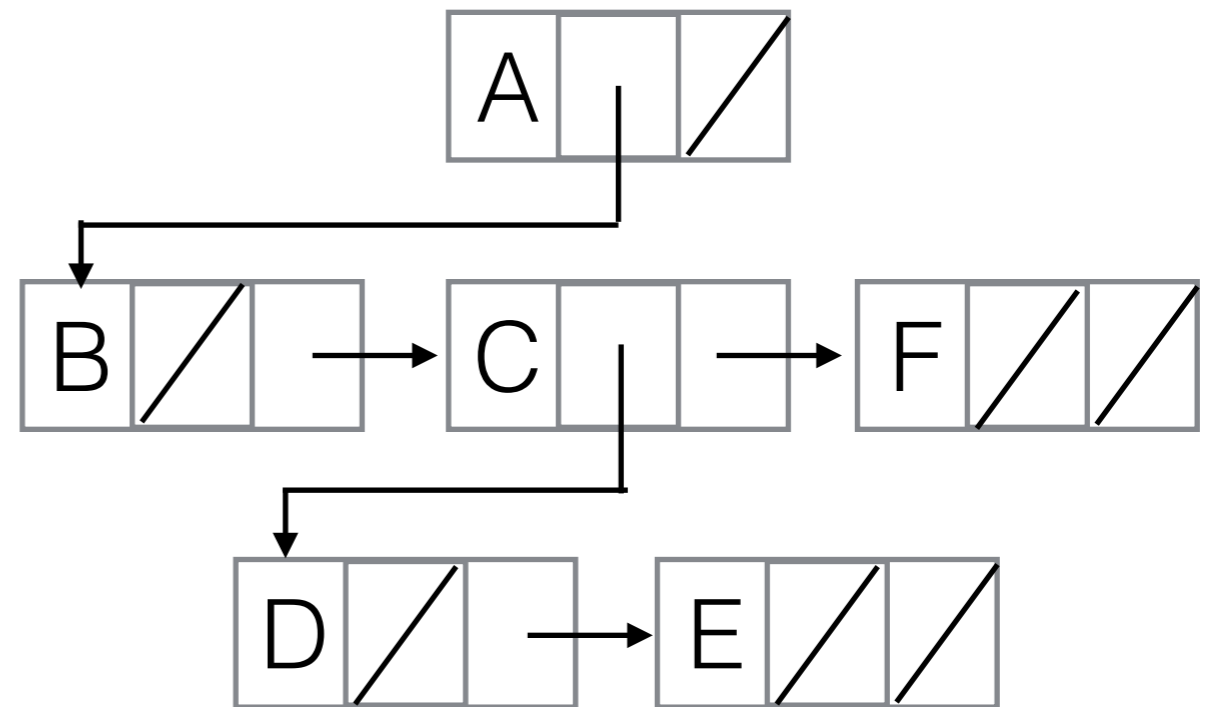
Traversal

- Using recursion
- Pre-Order: **VL to R**
 - A B C D E F
- In-Order: **LV to R**
 - B A D C E F
- Post-Order: **L to RV**
 - B D E C F A
- Some traversals more useful than others depending on situation
 - recall BSP



"General Tree" Implementation

- Q. Remember how to implement a tree using arrays?
- Any number of children
- Using pointers a tree is similar to a linked list
 - leftmost child pointer
 - right sibling pointer
- A **limited tree** with fixed number of children would be easier



Binary Trees

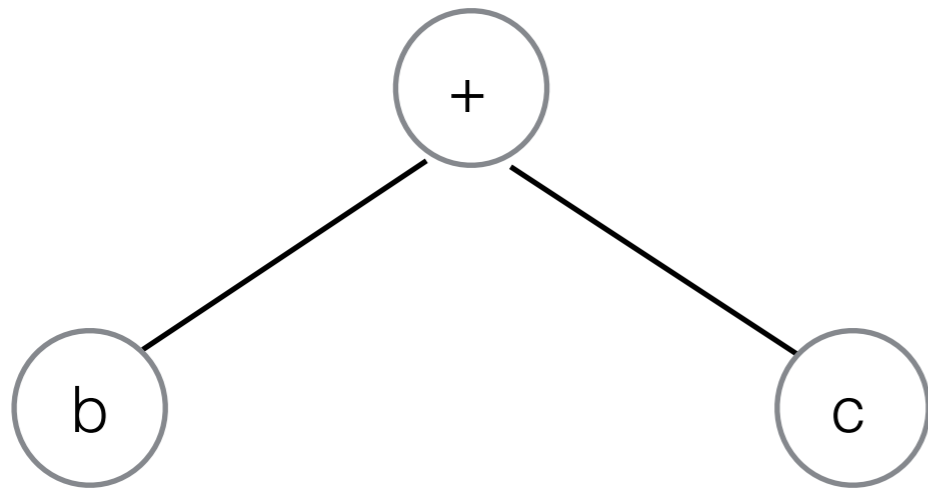
- 2 branches or less per node
- traversals only have 2 subtrees
- VLR (pre order), LVR (in-order), LRV (post-order)
- writing a *function* to add a new node is a little tricky (as with linked lists)
- tutorial?

```
struct Tree_Node;
struct Tree_Node {
    char data;
    Tree_Node *left, *right;
};

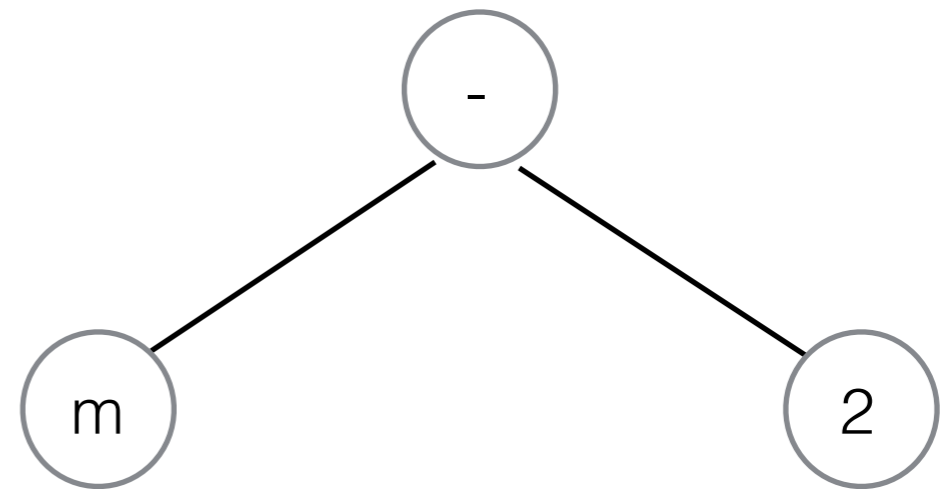
Tree_Node *root = (Tree_Node*)malloc(
    sizeof(Tree_Node));
root->data = 'A';
root->left = NULL;
root->right = NULL;
```

Arithmetic Trees

- construct tree to represent arithmetic expression
 $a * (b + c) / (m - 2)$
- set up sub-tree for each set of brackets



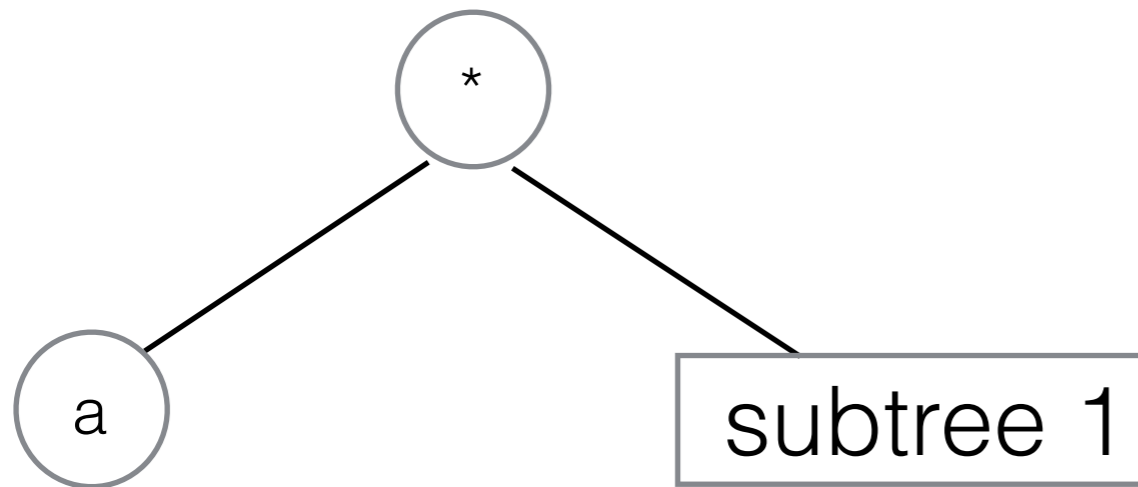
subtree 1



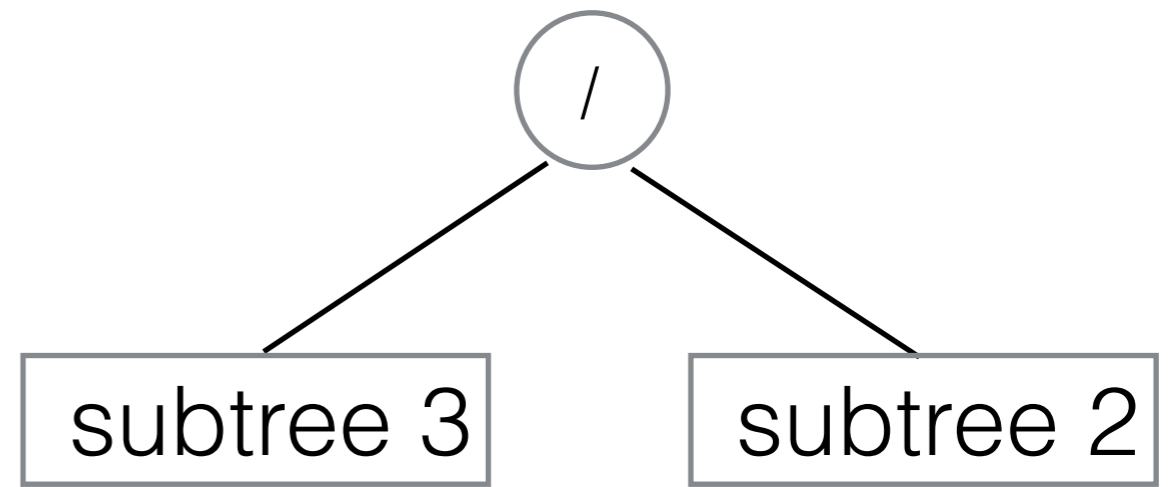
subtree 2

Arithmetic Trees

- $\mathbf{a * (b + c) / (m - 2)}$



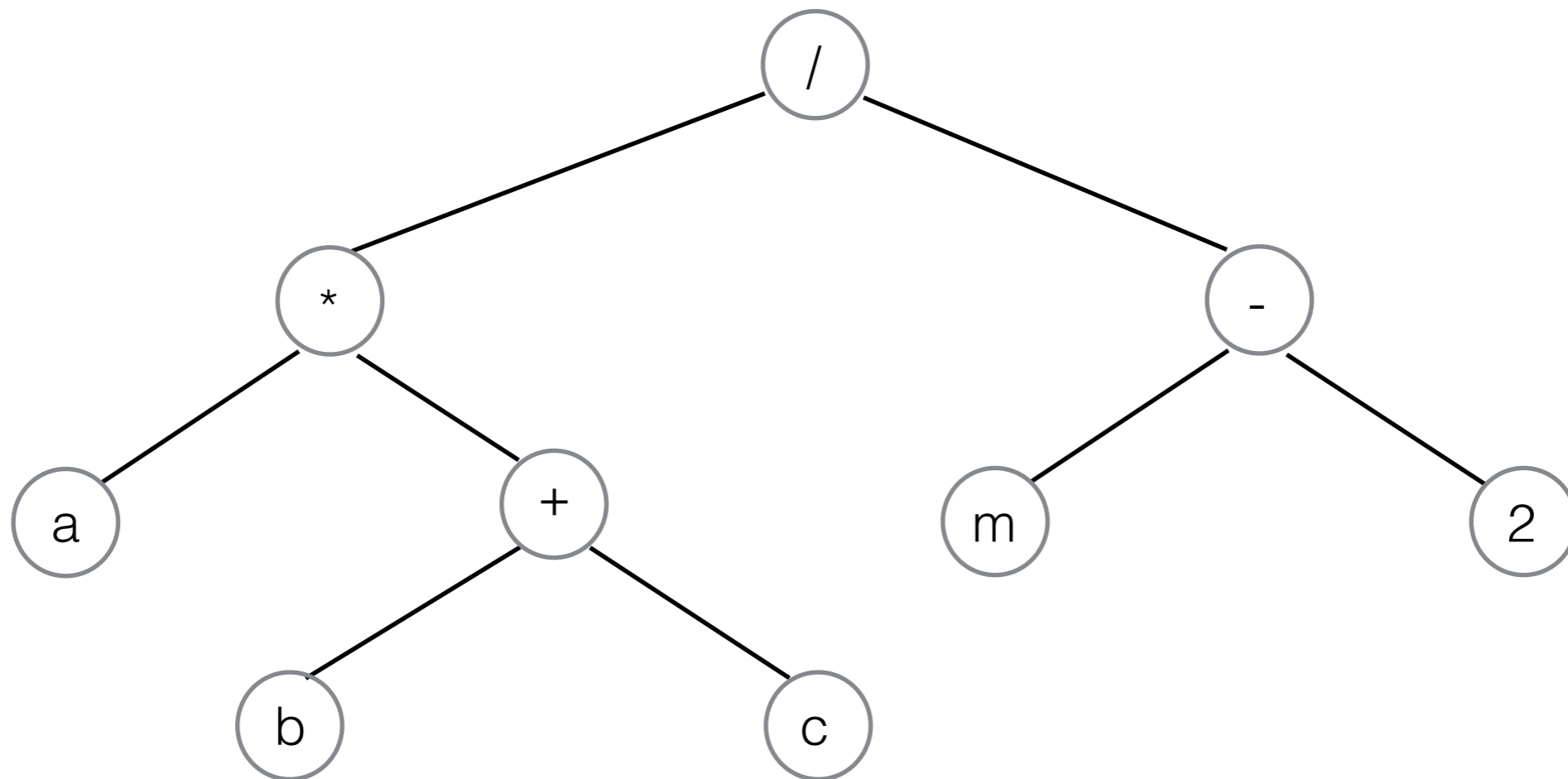
subtree 3



finally

Arithmetic Trees

- **Q.** in-order traversal generates: ... ? (that's L,V,R)



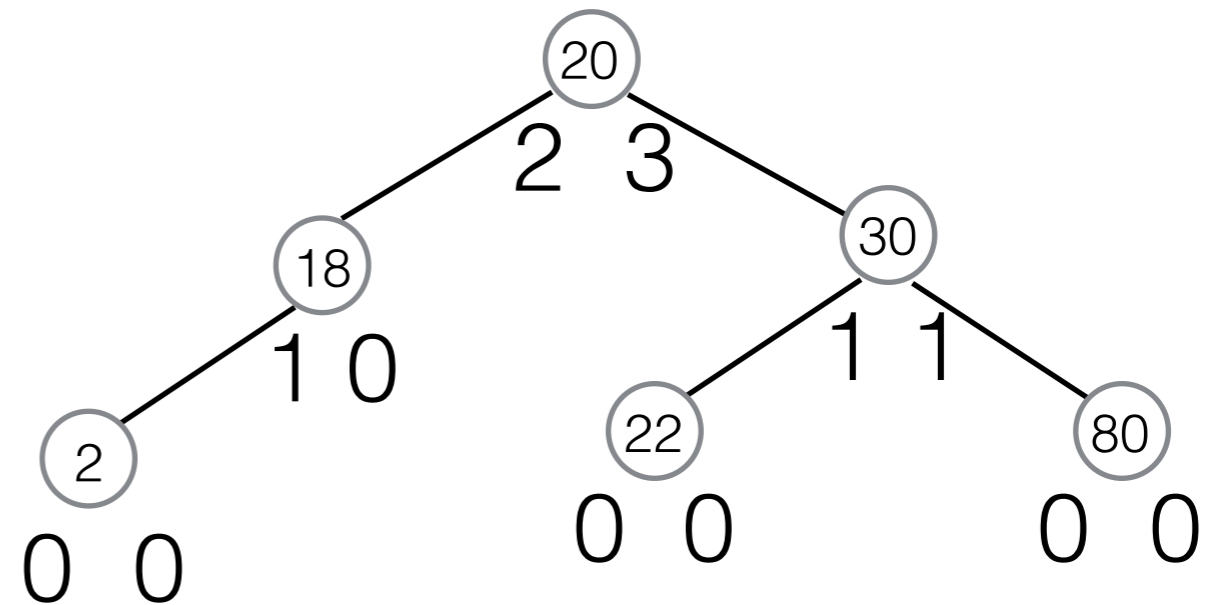
see also: Reverse Polish Notation

Binary Search Tree (BST)

- uses a **binary tree**
- data stored in any node is unique
- any data in **left** subtree is **less than root**
- any data in **right** subtree is **greater than root**
- left and right subtrees are also binary trees

Balance

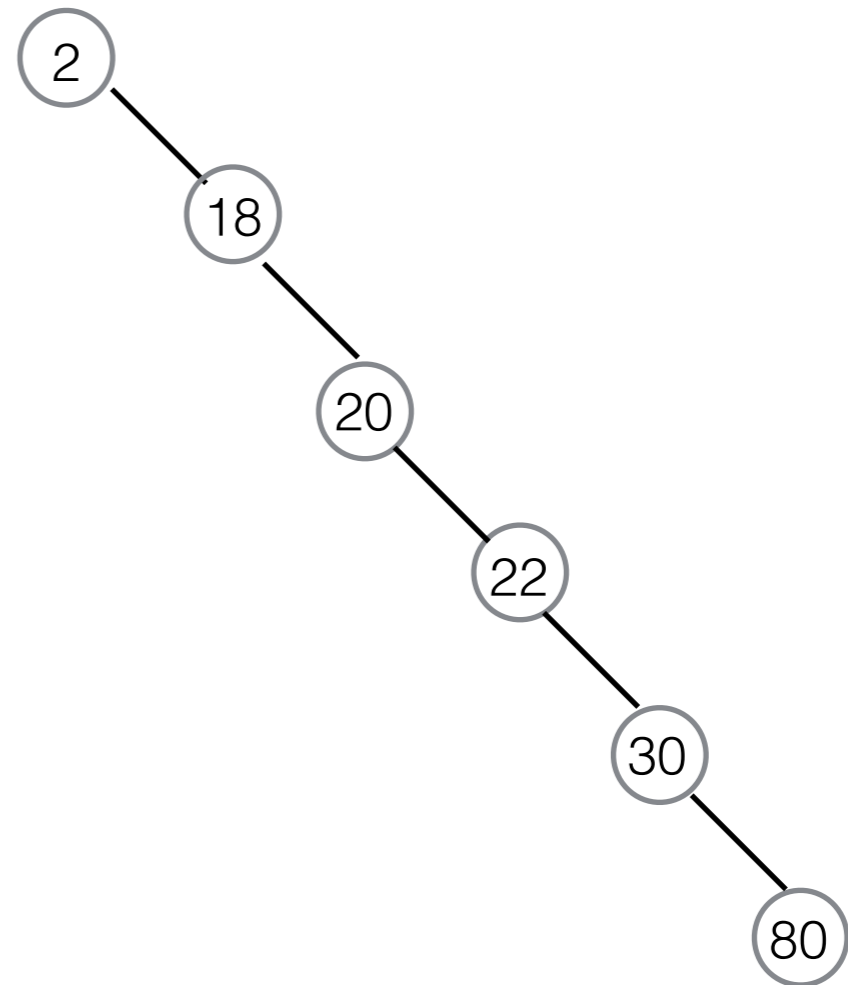
- A binary tree is perfectly **balanced** if
 - total nodes in left and right trees differs by max 1
 - levels in left and right trees differs by max 1



- To find if 80 is in this tree:
 - compare: 20, 30, 80

Balance

- Not balanced, but still a binary search tree:
 - 6 comparisons needed to find 80
- Search on BST
 - worst case **$O(n)$**
 - average case **$O(\log n)$**



Balance

- A tree will be balanced if we *insert* values in particular order
 - 20, 18, 30, 2, 22, 80 - first tree
 - 2, 18, 20, 22, 30, 80 - second tree
- If we **sort** the data into a list or array first we can create a perfectly balanced tree:
 - 2, 26, 30, 34, 56, 60, **65**, 70, 80, 94, 96, 98, 99
- *Then* we can choose the **mid value** - 65 - insert that first, split into a left and right list - choose mids of those, and so on.
 - so insertion order will be 65, 30, 2, 26, 56, 34, 60...