

Sorting Algorithms

Part 2

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Programming

- More live coding (maybe less theory time)
- Schedule fixed practice for most important skills
 - 30mins/day competent habit, 1hr+ competitive
 - I used to sit with flatmates semi-watching TV with the laptop
 - Code all the elementary sorts
 - Ask questions / discuss / ask for help / feedback
- Join DUCSS and Netsoc - they have some good talks/workshops

Previously...

- Charts for visualising sorting happening
- Terms: *file, record, key, stable, indirect sort*
- Selection sort - select smallest item from unsorted part (right) of array into current position
- Insertion sort - insert next item into correct position inside sorted part (left) of array (pickup pile and hand of sorted cards)

Bubble Sort Algorithm

- sorted = false
- while (sorted == false)
 - sorted = true
 - loop over data
 - if (next < current)
 - swap(current, next)
 - sorted = false

Bubble Sort

- Time and space complexity? Worst, average, **best**?
- Advantages:
 - Code is simple
 - Can stop early if numbers already sorted
 - No other sorting algorithm does this
 - Can do one run to check before calling complex sorting algorithm
 - "Stable"
- *Sedgewick has a different implementation of Bubble Sort
- Computer scientists have very negative things to say about Bubble Sort's worst case performance vs Insertion Sort.

Summary - Elementary Sorting Algorithms

- Very simple to implement. Also interchangeable. Some useful properties.
- $O(n^2)$ worst case time
 - May not play well with cache - try them with a timer
- $O(1)$ auxiliary memory (1 variable for swapping)
- Stable
- Code all of these yourself as exercise

merge() 2 sorted lists

- I have 2 sorted files (or arrays) A and B - merge them into a new output array
- Create 3 iterators (counters), one per array
 - `int a_index = 0, b_index = 0, output_index = 0;`
- Compare the value at each index, find the smallest
 - copy value to a new array
 - increment counter of the list you copied from
- Add any left-overs from A and B to output

Working Down the Page

List A	List B	Output List	<i>could track counters too</i>
3 4 12	1 10 23		
3 4 12	1 10 23	: 1	
3 4 12	1 10 23	: 1, 3	
3 4 12	1 10 23	: 1, 3, 4	
...	...	: ...	
		: 1, 3, 4, 10, 12, 23	

Merging

- Fairly simple
- $O(N)$
- Requires auxiliary memory - how much?
- Should I code this now? Might take a while - error prone.

Merge Sort Algorithm

1. Cut array of keys in half
2. Sort left half (recursively)
3. Sort right half (recursively)
4. Merge the two sorted lists

Merge Sort

- Merging two sorted lists is $O(n)$
- Bisecting the sort space is $O(\log(n))$
- So the whole sort is $O(n * \log(n))$
- Faster than our $O(n^2)$ elementary sorting algorithms
- More complex to implement
- Auxiliary memory use? $O(\dots)$

coding merge() for merge_sort()

- took me over an hour to code correctly
 - always print output and know what result *should* be
- made lots of mistakes and had to use the debugger
 - mixing up index variables
 - using `<` instead of `<=`
 - had to create a temp array inside `merge()` to avoid overwriting original data
- simplified my code after looking at others' code
- replacing recursion with loops would be better still

If we have time..

- Coding `merge()` and `merge_sort()` — might take too long for lecture - maybe in tutorial?
- One I prepared earlier follows (and link to GitHub in Discussion Board)
- Next: Quicksort, sorting and coding exercise, 2nd lab for assignment.

```

// first and last are the range of the output list, inclusive
// first half is left list, second half of this is the right list
void merge( int first, int last, int *array ) {
    // make a temporary working array so we don't overwrite our data
    // as we are reading it
    // alloca is dynamic _stack_ memory - freed at function close
    // you could do this with another sort of array or memory
    int* result = alloca(sizeof(int) * (last - first));

    int mid_index = ( first + last ) / 2;
    int left_index = first, right_index = mid_index + 1, output_index = first;

    // compare the lists until one list runs out of list
    while ( left_index <= mid_index && right_index <= last ) {
        if ( array[left_index] < array[right_index] ) {
            result[output_index++] = array[left_index++];
        } else {
            result[output_index++] = array[right_index++];
        }
    }

    // copy any leftovers from either list into output
    // you can probably simplify these into the other loop
    // if you're smarter than me
    while ( left_index <= mid_index ) {
        result[output_index++] = array[left_index++];
    }
    while ( right_index <= last ) {
        result[output_index++] = array[right_index++];
    }

    // copy into original array
    for (int i = first; i <= last; i++) {
        array[i] = result[i];
    }
}

```

```

// declare here so i can recursively call self
void merge_sort( int first_index, int last_index, int* data );

void merge_sort( int first_index, int last_index, int* data ) {
    // break recursion when counters meet in the middle
    if ( first_index >= last_index ) {
        return;
    }
    int mid_index = ( first_index + last_index ) / 2;

    // NB: replacing recursion with loops is usually more efficient
    merge_sort( first_index, mid_index, data );
    merge_sort( mid_index + 1, last_index, data );
    merge( first_index, last_index, data );
}

int main() {
    // create 2 input lists and space for one output list
    int data[] = { 3, 4, 12, 1, 10, 23 }; // initialiser list for array
    giving constant values

    // sort with bisections, recursively, from indices 0 to 5, inclusive
    merge_sort( 0, 5, data );

    for ( int i = 0; i < 6; i++ ) {
        printf( "%i ", data[i] );
    }
    printf( "\n" );

    return 0;
}

```